## CS 465: Language, Logic and Computation–Lambda Calculus Homework

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This homework is based on the material presented in lecture; you should refer to on-line hand-outs and your own lecture notes for definitions, etc. Please ask for help if you have any questions.

## 1. Syntactic abbreviations

Convert the following lambda terms into fully parenthesized form, with all multi-variable abstractions expanded. Let's agree that "fully parenthesized" will mean that all applications and abstractions should be parenthesized, except for the outermost term (i.e., at top level).

- λxyz. f x (f y y) (f z)
- $\lambda xy. f(\lambda x. x)(\lambda gf. gf x)$

Perform the "opposite operation" on the following terms; i.e., use standard conventions on parenthesization and multiple abstractions to write them in a *minimal* form.

- λa. (λb. (a ((b b) a))
- $(((\lambda p. (\lambda x. (p x)) (p a)) (\lambda q.q)) (\lambda y.b)) a$

## 2. Variables and binding

Convert the following terms so that no free or bound variables clash, i.e., so that all variables are distinct (respect the existing variable bindings: you may change the variables, but preserve the meaning).

- (λx. (λxy. x (y x)) y) (λy. (λy. y x) y x)
- b a  $(\lambda ab. a (\lambda a. b) b)$  a  $(\lambda b. b a)$

## 3. Substitution and reduction

Reduce the following lambda and combinator terms to normal form, using normal-order reduction, or argue that they have no normal form. Show intermediate  $\beta$ -reduction steps and, if necessary, variable renamings to avoid capture.

- (λfgx. f (g x)) (λy.y) (λfx. f (f x)) x
- (λfx. f (f x)) (λy. c (y y)) b
- SII(SII)